

# Recent Advances in IPv6 Security

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# About...

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- Security researcher and consultant at SI6 Networks
- Have worked on security assessment on communications protocols for:
  - UK NISCC (National Infrastructure Security Co-ordination Centre)
  - UK CPNI (Centre for the Protection of National Infrastructure)
- Active participant at the IETF (Internet Engineering Task Force)
- More information available at: <http://www.gont.com.ar>

# Agenda

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- Disclaimer
- Motivation for this presentation
- Recent Advances in IPv6 Security
  - IPv6 Addressing
  - IPv6 Fragmentation & Reassembly
  - IPv6 First Hop Security
  - IPv6 Firewalling
  - Mitigation to some Denial of Service attacks
- Conclusions
- Questions and Answers

# Disclaimer

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- This talk assumes:
  - You know the basics of IPv4 security
  - You know the basics about IPv6 security
  - (i.e. I'm not doing an “IPv6 primer” in this presentation, sorry)
- Much of this is “work in progress” → your input is welcome!
- No “0-days”, sorry.

# Motivation for this presentation

# Motivation for this presentation

- Sooner or later you will need to deploy IPv6
  - In fact, you have (at least) partially deployed it, already
- IPv6 represents a number of challenges: What can we do about them?

## Option #1

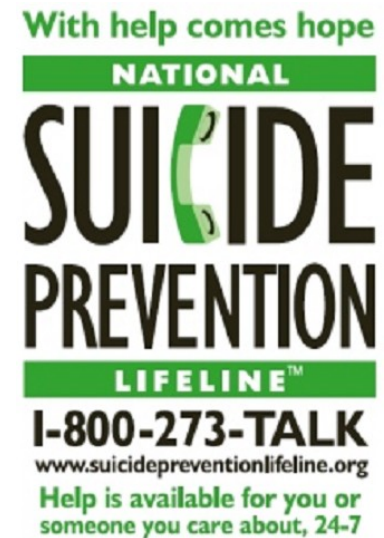


## Option #2



Suicide is always an option.

## Option #3



# Motivation for this presentation (II)

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- We have been doing a fair share of IPv6 security research
  - Identification of problems
  - Proposals to mitigate those problems
- Part of our research has been taken to the IETF
- This talk is about our ongoing work to improve IPv6 security

# Advances in IPv6 Addressing



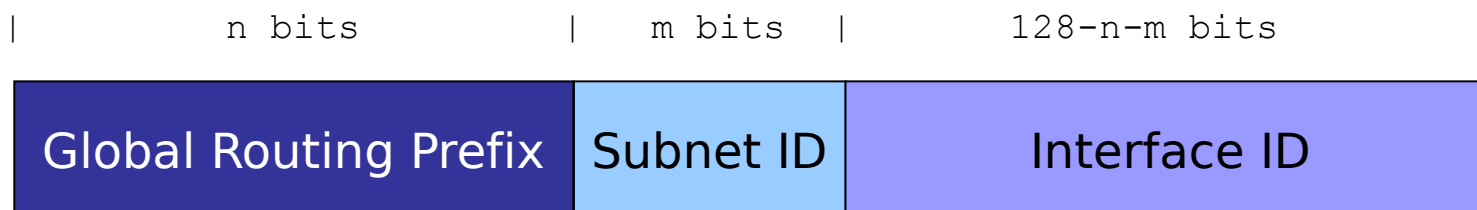
# Brief overview

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- Main driver for IPv6 deployment
- Employs 128-bit addresses
- Address semantics similar to those of IPv4:
  - Addresses are aggregated into “prefixes”
  - Several address types
  - Several address scopes
- Each interface typically employs more than one address, of different type/scope:
  - One link-local unicast address
  - One or more global unicast addresses
  - etc.

# Global Unicast Addresses

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- The “Interface ID” is typically 64-bit long
- Can be selected with different criteria:
  - Modified EUI-64 Identifiers
  - Privacy addresses
  - Manually configured
  - As specified by transition/co-existence technologies

# IPv6 Addressing

## Implications on remote address scanning attacks

# IPv6 host scanning attacks

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“Thanks to the increased IPv6 address space, IPv6 host scanning attacks are unfeasible. Scanning a /64 would take 500.000.000 years”

– Urban legend

**Is the search space for a /64 really  $2^{64}$  addresses?**

# IPv6 addresses in the real world

- Malone measured (\*) the address generation policy of hosts and routers in real networks

Address type	Percentage
SLAAC	50%
IPv4-based	20%
Teredo	10%
Low-byte	8%
Privacy	6%
Wordy	<1%
Others	<1%

Hosts

Address type	Percentage
Low-byte	70%
IPv4-based	5%
SLAAC	1%
Wordy	<1%
Privacy	<1%
Teredo	<1%
Others	<1%

Routers

Malone, D., "Observations of IPv6 Addresses", Passive and Active Measurement Conference (PAM 2008, LNCS 4979), April 2008, <<http://www.maths.tcd.ie/~dwmalone/p/addr-pam08.pdf>>.

# IPv6 addresses embedding IEEE IDs



- In practice, the search space is at most  $\sim 2^{24}$  bits – **feasible!**
- The low-order 24-bits are not necessarily random:
  - An organization buys a large number of boxes
  - In that case, MAC addresses are usually consecutive
  - Consecutive MAC addresses are generally in use in geographically-close locations

# IPv6 addresses embedding IEEE IDs (II)

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- Virtualization technologies present an interesting case
- Virtual Box employs OUI 08:00:27 (search space:  $\sim 2^{24}$ )
- VMWare ESX employs:
  - Automatic MACs: OUI 00:05:59, and next 16 bits copied from the low order 16 bits of the host's IPv4 address (search space:  $\sim 2^8$ )
  - Manually-configured MACs: OUI 00:50:56 and the rest in the range 0x000000-0x3ffff (search space:  $\sim 2^{22}$ )

# IPv6 addresses embedding IPv4 addr.

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- They simply embed an IPv4 address in the IID
  - e.g.: 2000:db8::192.168.0.1
- Search space: same as the IPv4 search space



# IPv6 “low-byte” addresses

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- The IID is set to all-zeros, except for the last byte
  - e.g.: 2000:db8::1
  - There are other variants
- Search space: usually  $2^8$  or  $2^{16}$

# Industry mitigations for scanning attacks

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- Microsoft replaced the MAC-address-based identifiers with (non-standard) randomized IIDs
  - Essentially RFC 4941, but they don't vary over time
- Certainly better than MAC-address-based IIDs, but still not “good enough”
- They mitigate host-scanning, but **not** host tracking – constant IIDs are still present!

# Thoughts on remote scanning attacks

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- IPv6 host scanning attacks are **feasible**, but typically harder than in IPv4
- They require more “intelligence” on the side of the attacker
- It is **possible** to make them infeasible
- It is likely that many other scanning strategies/techniques will be explored

# IPv6 Addressing

## Implications on privacy

# Problem statement

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- Modified EUI-64 IIDs are constant for each interface
- As the host moves, the prefix changes, but the IID doesn't
  - the 64-bit IID results in a super-cookie!
- This introduces a problem not present in IPv4: **host-tracking**
- Example:
  - In net #1, host configures address: 2001:db8:1::1111:2222:3333:4444
  - In net #2, host configures address: 2001:db8:2::1111:2222:3333:4444
  - The IID “1111:2222:3333:4444” leaks out host “identity”.

# “Mitigation” to host-tracking

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- RFC 4941: privacy/temporary addresses
  - Random IIDs that change over time
  - Generated **in addition** to traditional SLAAC addresses
  - Traditional addresses used for server-like communications, temporary addresses for client-like communications
- Operational problems:
  - Difficult to manage!
- Security problems:
  - They mitigate host-tracking **only partially**
  - They **do not** mitigate address-scanning attacks

# **IPv6 addressing**

## **Mitigating scanning and privacy issues**

# Auto-configuration address/ID types

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	Stable	Temporary
Predictable	IEEE ID-derived	None
Unpredictable	<b>NONE</b>	RFC 4941

- We lack stable privacy-enhanced IPv6 addresses
  - Used to replace IEEE ID-derived addresses
  - Pretty much orthogonal to privacy addresses
  - Probably “good enough” in most cases even without RFC 4941



# Stable privacy-enhanced addresses

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- draft-ietf-6man-stable-privacy-addresses proposes to generate Interface IDs as:

$F(\text{Prefix}, \text{Interface\_Index}, \text{Network\_ID}, \text{Secret\_Key})$

- Where:
  - $F()$  is a PRF (e.g., a hash function)
  - Prefix SLAAC or link-local prefix
  - Interface\_Index is the (internal) small number that identifies the interface
  - Network\_ID could be e.g. the SSID of a wireless network
  - Secret\_Key is unknown to the attacker (and randomly generated by default)

# Stable privacy-enhanced addresses (II)

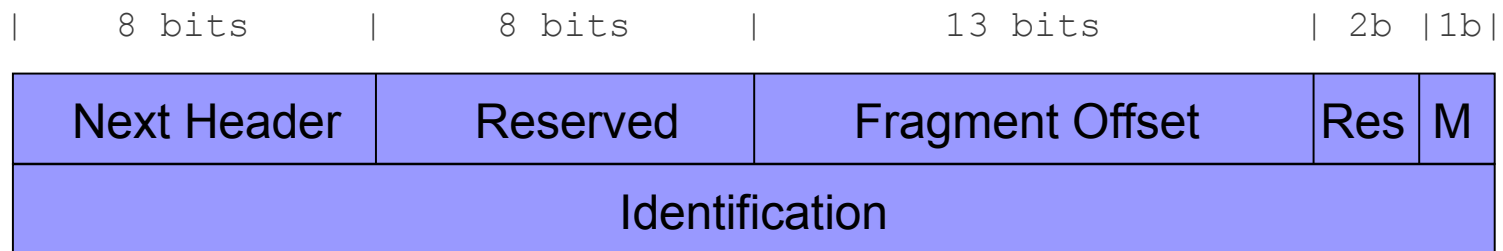
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- As a host moves:
  - Prefix and Network\_ID change from one network to another
  - But they remain constant within each network
  - F() varies across networks, but remains constant within each network
- This results in addresses that:
  - Are stable within the same subnet
  - Have different Interface-IDs when moving across networks
  - For the most part, they have “the best of both worlds”
- Document already accepted as a 6man wg item

# IPv6 Fragmentation and Reassembly

# IPv6 fragmentation

- IPv6 fragmentation performed only by hosts (never by routers)
- Fragmentation support implemented in “Fragmentation Header”
- Fragmentation Header syntax:



# Fragment Identification

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- Security Implications of predictable Fragment IDs well-known from the IPv4 world
  - idle-scanning, DoS attacks, etc.
- Amount of fragmented traffic will probably increase as a result of:
  - Larger addresses
  - DNSSEC
- But no worries, since we learned the lesson from the IPv4 world... – **right?**

# Fragment ID generation policies

Operating System	Algorithm
FreeBSD 9.0	Randomized
NetBSD 5.1	Randomized
OpenBSD-current	Randomized (based on SKIPJACK)
Linux 3.0.0-15	<b>Predictable</b> (GC init. to 0, incr. by +1)
Linux-current	Unpredictable (PDC init. to random value)
Solaris 10	<b>Predictable</b> (PDC, init. to 0)
Windows 7 Home Prem.	<b>Predictable</b> (GC, init. to 0, incr. by +2)

GC: Global Counter      PDC: Per-Destination Counter

At least Solaris and Linux patched in response to our IETF I-D – more patches expected!

# Fixing predictable Fragment IDs

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- draft-gont-6man-predictable-fragment-id:
  - Discussed the security implications of predictable Fragment ID
  - Proposes a number of algorithms to generate the Fragment ID
- Ongoing work at the 6man wg
  - Has not yet been adopted by the 6man working group

# IPv6 Fragment Reassembly

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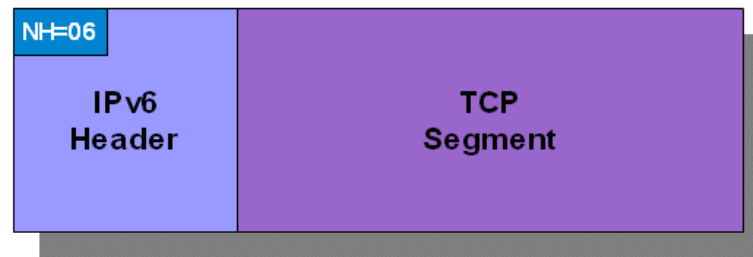
- Security implications of overlapping fragments well-known (think Ptacek & Newsham, etc,)
- Nonsensical for IPv6, but originally allowed in the specs
- Different implementations allow them, with different results
- RFC 5722 updated the specs, forbidding overlapping fragments
- Most current implementations reflect the updated standard
- See <http://blog.si6networks.com>



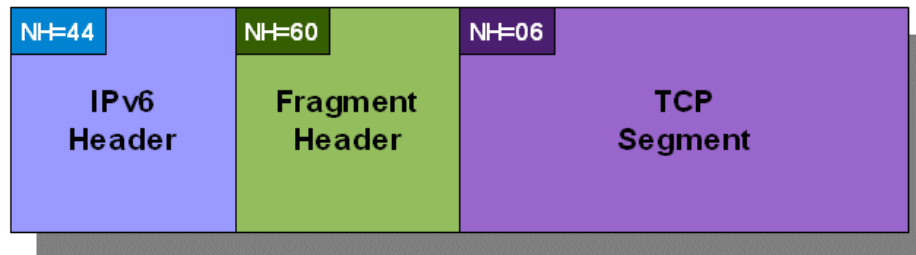
# IPv6 “atomic” fragments

- ICMPv6 PTB < 1280 triggers inclusion of a FH in all packets to that destination (not actual fragmentation)
- Result: IPv6 atomic fragments (Frag. Offset=0, More Frag.=0)

**Original packet**



**Atomic fragment**



# Issues with IPv6 atomic fragments

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- Some implementations mix “atomic fragments” with queued fragments
- Atomic fragments thus become subject of IPv6 fragmentation attacks
- How to leverage this issue:
  - Trigger atomic fragments with ICMPv6 PTB messages
  - Now perform IPv6 fragmentation-based attacks

# Mitigating issues with atomic fragments

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- draft-ietf-6man-ipv6-atomic-fragments fixes the problem:
  - IPv6 atomic fragments required to be processed as non-fragmented traffic
- Document has passed WGLC
  - Should be published as an RFC this year

# Handling of IPv6 atomic fragments

Operating System	Atomic Frag. Support	Improved processing
FreeBSD 8.0	No	No
FreeBSD 8.2	Yes	No
FreeBSD 9.0	Yes	No
Linux 3.0.0-15	Yes	Yes
NetBSD 5.1	No	No
OpenBSD-current	Yes	Yes
Solaris 11	Yes	Yes
Windows Vista (build 6000)	Yes	No
Windows 7 Home Premium	Yes	No

At least OpenBSD patched in response to our IETF I-D – more patches expected!

# IPv6 First Hop Security

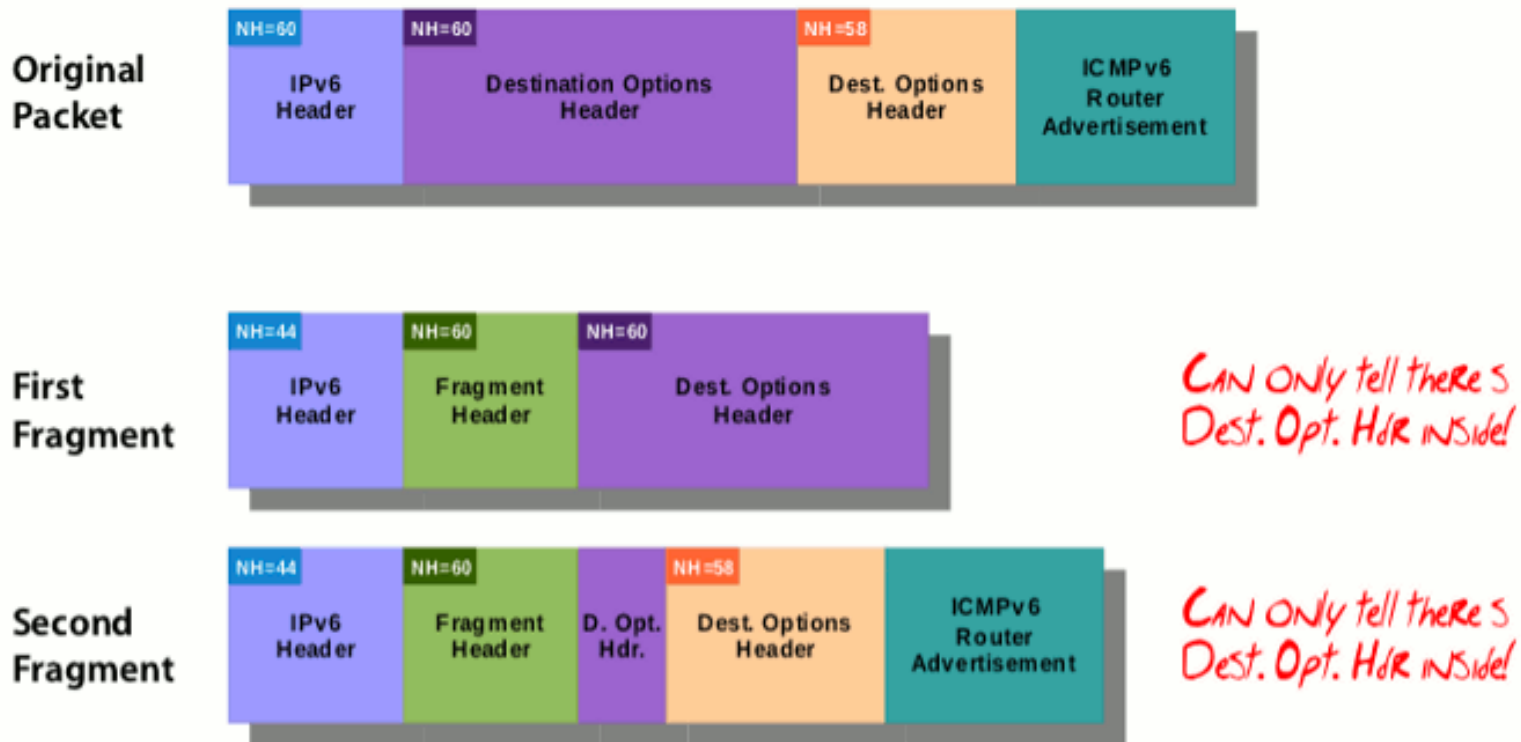
# IPv6 First Hop Security

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- Security mechanisms/policies employed/enforced at the first hop (local network)
- Fundamental problem: lack of feature-parity with IPv4
  - arpwatc-like Neighbor Discovery monitoring virtually impossible
  - DHCP-snooping-like RA blocking trivial to circumvent

# IPv6 First Hop Security (II)

- Fundamental problem: complexity of traffic to be “processed at layer-2”
- Example:



# Bringing “sanity” to ND traffic

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- draft-ietf-6man-nd-extension-headers forbids use of fragmentation with Neighbor Discovery
  - It makes ND monitoring feasible
  - Turns out it is vital for SEND (or SEND could be DoS'ed with fragments)
- Work in progress:
  - Has been adopted as a 6man wg item
  - Should be published as an RFC “shortly”



# RA-Guard

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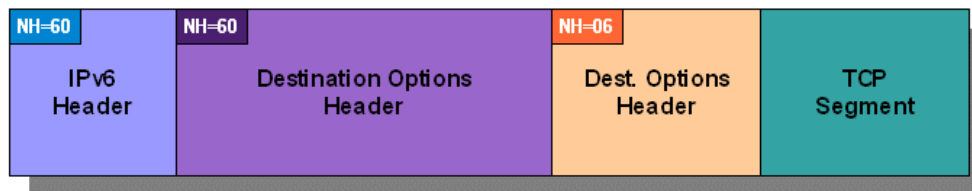
- Meant to block RA packets on “unauthorized” switch ports
- Existing implementations trivial to circumvent
- draft-ietf-v6ops-ra-guard-implementation contains:
  - Discussion of RA-Guard evasion techniques
  - Advice to filter RAs, while avoiding false positives
- Can only be evaded with overlapping fragments
  - But most current OSes forbid them
  - And anyway there's nothing we can do about this :-)
- Should be published as an RFC this year.

# IPv6 firewalling

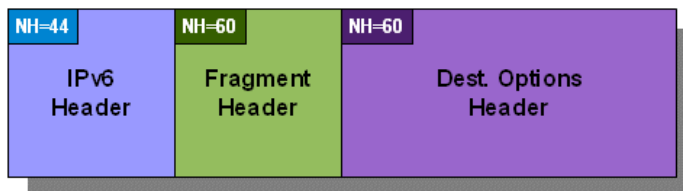
# Problem statement

- Specs-wise, state-less IPv6 packet filtering is impossible:
  - The IPv6 header chain can span multiple fragments
  - This makes state-less firewalling impossible

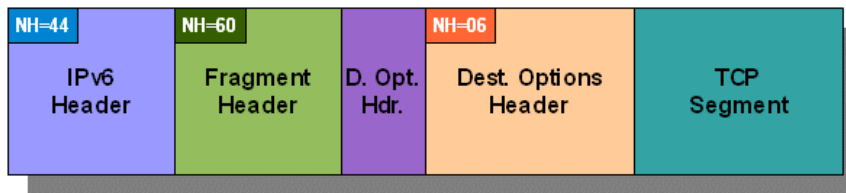
Original packet



First fragment



Second fragment



# First step away from “insanity”

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- draft-ietf-6man-oversized-header-chain fixes this problem:
  - The entire IPv6 header chain must be contained in the first fragment
  - i.e. packets with header chains that span more than one fragment may be blocked – don't send them!
- Work in progress:
  - Already adopted by the 6man WG
  - Should be published as an RFC “shortly”
- There's an insanely large amount of work to be done in the area of IPv6 firewalling

# IPv6 implications on IPv4 networks

# VPN leakages

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- Typical scenario:
  - You connect to an insecure network
  - You establish a VPN with your home/office
  - **Your VPN software does not support IPv6**
- Trivial to trigger a VPN leakage
  - Spoof RA's or DHCPv6-server packets, to set the recursive DNS server
  - Simply trigger IPv6 connectivity, such that dual-stacked hosts leak out
  - Even legitimate dual-stacked networks may trigger it
- As always, deemed as “already known” by some
  - Yet most VPN clients are vulnerable, and nobody did anything about it

# Tools

# IPv6 security tools

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- For ages, THC's IPv6 attack suite (<http://www.thc.org>) was the only IPv6 security tools publicly available
- We've produced “SI6 Networks IPv6 toolkit”
  - A brand-new security assessment/trouble-shooting toolkit
  - Runs on Linux, \*BSD, and Mac OS
- Available at: <http://www.si6networks.com/tools/ipv6toolkit>
  - GIT repository at: <https://github.com/fgont/ipv6-toolkit.git>



# SI6 Networks' IPv6 toolkit

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- scan6: An IPv6 address scanner
- frag6: Play with IPv6 fragments
- tcp6: Play with IPv6-based TCP segments
- ns6: Play with Neighbor Solicitation messages
- na6: Play with Neighbor Advertisement messages
- rs6: Play with Router Solicitation messages
- ra6: Play with Router Advertisement messages

# SI6 Networks' IPv6 toolkit (II)

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- rd6: Play with Redirect messages
- icmp6: Play with ICMPv6 error messages
- ni6: Play with Node Information messages
- flow6: Play with the IPv6 Flow Label
- jumbo6: Play with IPv6 Jumbograms

# Some conclusions

# Some conclusions

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- Many IPv4 vulnerabilities have been re-implemented in IPv6
  - We just didn't learn the lesson from IPv4, or,
  - Different people working in IPv6 than working in IPv4, or,
  - The specs could make implementation more straightforward, or,
  - **All of the above? :-)**
- Still lots of work to be done in IPv6 security
  - We all know that there is room for improvements
  - **We need IPv6, and should work to improve it**

# Questions?

# Thanks!

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IPv6 Hackers mailing-list

<http://www.si6networks.com/community/>



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